



NO-REGRET LEARNING IN GAMES

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Outline

Background

N-player games

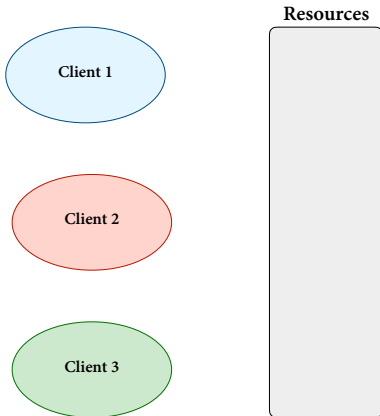
No-regret learning in games

Learning with limited feedback



The Kelly auction

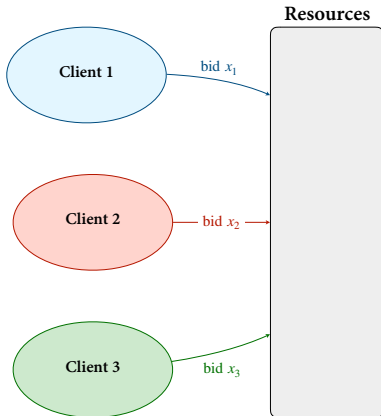
Proportionally fair allocation of resources to different clients [Kelly, 1998]:





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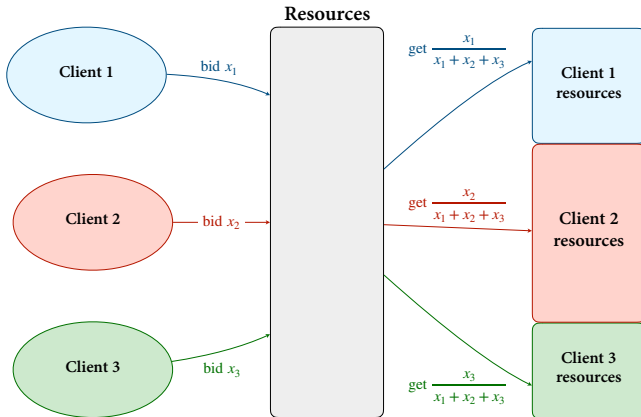
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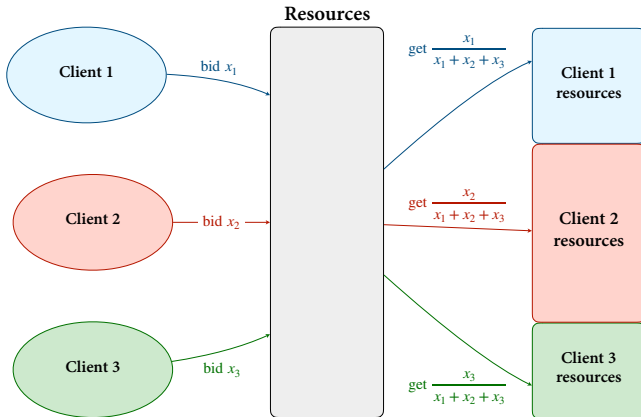
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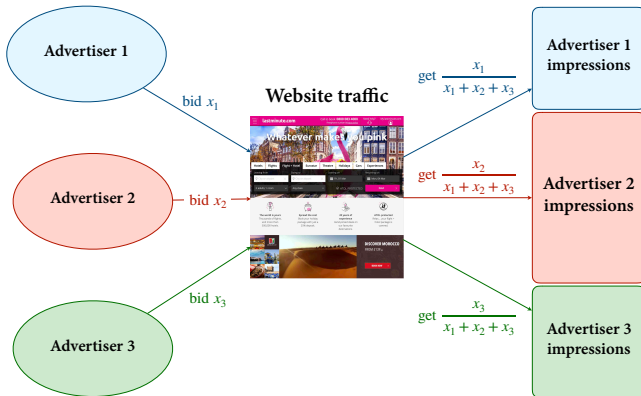


Resources could be processor cores, bandwidth



The Kelly auction

Proportionally fair allocation of resources to different clients [Kelly, 1998]:



Resources could be processor cores, bandwidth , or even **anonymous web traffic**



Online decision processes

Agents called to take repeated decisions with minimal information:

repeat

At each epoch $t = 1, 2, \dots$

Choose **action** X_t

Get **payoff** $u_t(X_t)$

until end



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At each epoch $t = 1, 2, \dots$

Choose **action** X_t

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Main question: *How to choose a “good” action at each epoch?*

- ▶ **Uncertain world:** no beliefs, feedback, knowledge of future, etc.
- ▶ **Obliviousness:** are payoffs affected by the agent’s previous actions?
- ▶ **Optimality:** what is “optimal” in this setting?



Regret minimization

Performance often quantified by the agent's **regret**

$$u_t(x) - u_t(X_t)$$



Regret minimization

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$$\sum_{t=1}^T [u_t(x) - u_t(X_t)]$$



Regret minimization

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$$\max_{x \in \mathcal{X}} \sum_{t=1}^T [u_t(x) - u_t(X_t)]$$



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$$\text{Reg}(T) = \max_{x \in \mathcal{X}} \sum_{t=1}^T [u_t(x) - u_t(X_t)]$$



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No regret: $\text{Reg}(T) = o(T)$

"The sequence of chosen actions is as good as the best fixed action in hindsight."



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Prolific literature:

- ▶ Economics [Hannan, Blackwell, Hart & Mas-Colell,...]
- ▶ Machine learning & computer science [Littlestone & Warmuth, Vovk,...]
- ▶ Online learning & optimization [Cesa-Bianchi & Lugosi, Zinkevich,...]



Multi-agent learning

- ▶ **Multiple** agents, individual objectives
- ▶ Payoffs determined by actions of **all** agents
- ▶ Agents receive payoffs, **adjust actions**, and the process repeats



Multi-agent learning

- ▶ **Multiple** agents, individual objectives
Example: place a bid in a repeated auction
- ▶ Payoffs determined by actions of **all** agents
Example: outcome of auction revealed
- ▶ Agents receive payoffs, **adjust actions**, and the process repeats
Example: change bid if unsatisfied



No-regret and equilibrium

The golden rule:

No-regret learning leads to equilibrium



No-regret and equilibrium

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No-regret learning leads to equilibrium*

* If it's ok to:



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X Assign positive weight only to strictly dominated strategies

[Viossat & Zapechelnyuk, 2013]



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X Be arbitrarily far from equilibrium infinitely often

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X ...



No-regret and equilibrium

When does no-regret learning converge to Nash equilibrium?



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N-player games

The game

- ▶ Finite set of *players* $i \in \mathcal{N} = \{1, \dots, N\}$
- ▶ Each player selects an *action* x_i from a compact convex set \mathcal{X}_i
- ▶ Reward of player i determined by *payoff function* $u_i: \mathcal{X} \equiv \prod_i \mathcal{X}_i \rightarrow \mathbb{R}$

Examples

- ▶ Finite games (mixed extensions)
- ▶ Power control/allocation problems
- ▶ Traffic routing
- ▶ Generative adversarial networks (two-player zero-sum games)
- ▶ Divisible good auctions (Kelly,...)
- ▶ Cournot oligopolies
- ▶ ...



Kelly auctions

The Kelly auction as an N -player game:

- ▶ **Players:** $i = 1, \dots, N$ [bidders]
- ▶ **Resources** $\mathcal{S} = \{1, \dots, S\}$ [websites]
- ▶ **Action spaces:** $\mathcal{X}_i = \{x_i \in \mathbb{R}_+^S : \sum_s x_{is} \leq b_i\}$ [b_i : budget of i -th bidder]
- ▶ **Resource allocation ratio:**

$$\rho_{is}(x) = \frac{q_s x_{is}}{c_{is} + \sum_{j \in \mathcal{N}} x_{js}}$$

[c_{is} : entry barrier]

- ▶ **Payoff functions:**

$$u_i(x) = \sum_{s \in \mathcal{S}} [g_i \rho_{is}(x) - x_i]$$

[utility from resources minus cost]



Nash equilibrium

Nash equilibrium

Action profile $x^* = (x_1^*, \dots, x_t^*) \in \mathcal{X}$ that is **unilaterally stable**

$$u_i(x_i^*; x_{-i}^*) \geq u_i(x_i; x_{-i}^*) \quad \text{for every player } i \in \mathcal{N} \text{ and every deviation } x_i \in \mathcal{X}_i$$



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Individual payoff gradients

$$V_i(x) = \nabla_{x_i} u_i(x_i; x_{-i})$$

Interpretation: direction of **individually** steepest payoff ascent



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Variational characterization

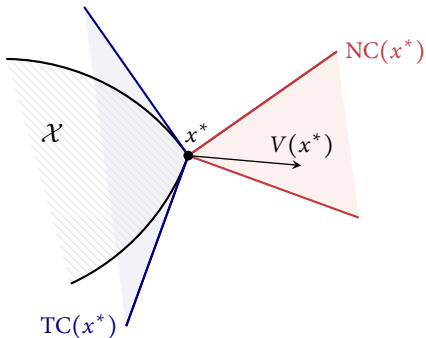
If x^* is a Nash equilibrium, then

$$\langle V_i(x^*), x_i - x_i^* \rangle \leq 0 \quad \text{for all } i \in \mathcal{N}, x_i \in \mathcal{X}_i$$

Intuition: $u_i(x_i; x_{-i}^*)$ decreasing along all rays emanating from x_i^*



Geometric interpretation



At Nash equilibrium, individual payoff gradients are outward-pointing



Monotonicity

A key assumption for games is **monotonicity**:

$$\langle V(x') - V(x), x' - x \rangle \leq 0 \quad \text{for all } x \in \mathcal{X} \quad (\text{MC})$$



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Equivalently: $H(x) \preceq 0$ where H is the game's **Hessian matrix**:

$$H_{ij}(x) = \frac{1}{2} \nabla_{x_j} \nabla_{x_j} u_i(x) + \frac{1}{2} (\nabla_{x_i} \nabla_{x_j} u_j(x))^\top$$

Interpretation: concavity for games



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Interpretation: concavity for games

Examples: Kelly auctions, Cournot oligopolies, routing, power control, ...

Close relatives:

- ▶ Stable games [Hofbauer & Sandholm, 2009]
- ▶ Contractive games [Sandholm, 2015];
- ▶ Dissipative [Sorin & Wan, 2016]



Monotonicity

Theorem (Rosen, 1965)

*If a game is strictly monotone, it admits a **unique Nash equilibrium**.*

[+ extensions to {...}-monotone games, generalized equilibrium problems,...]



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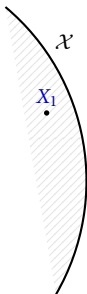


How to achieve no regret?

Take a gradient step and project:

[Zinkevich, ICML 2003]

$$X_{t+1} = \Pi(X_t + \gamma_t \nabla u_t(X_t)) \quad (\text{OGD})$$



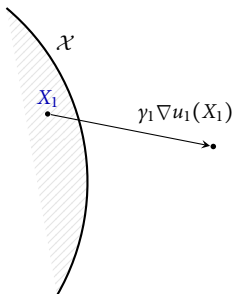


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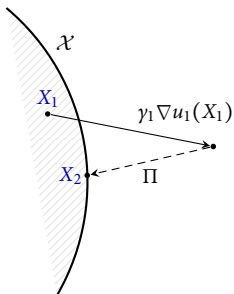


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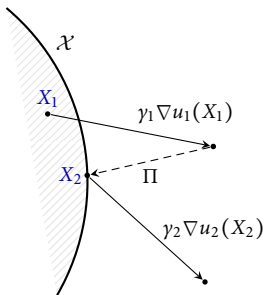


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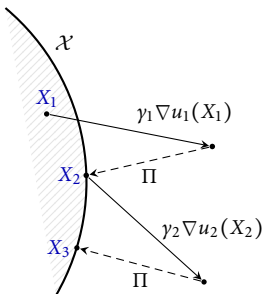


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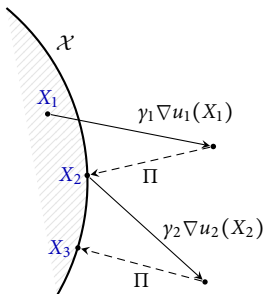


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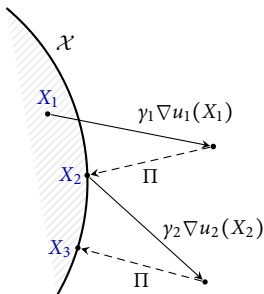


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...but what about convergence?



A dynamical systems viewpoint

Vector flow of V (simplest case: no constraints, smooth, etc.):

$$\frac{dX_i}{dt} = -V_i(X(t)) \quad (\text{GD})$$

Energy function:

$$E(x) = \frac{1}{2} \|x - x^*\|^2$$

Lyapunov property:

$$\frac{dE}{dt} = -\langle V(X(t)), X(t) - x^* \rangle \leq 0$$

Distance to solutions is (weakly) **decreasing** along trajectories of (GD)



Cycles

Roadblock: the energy might be a **constant of motion**

[Hofbauer et al, 2009]

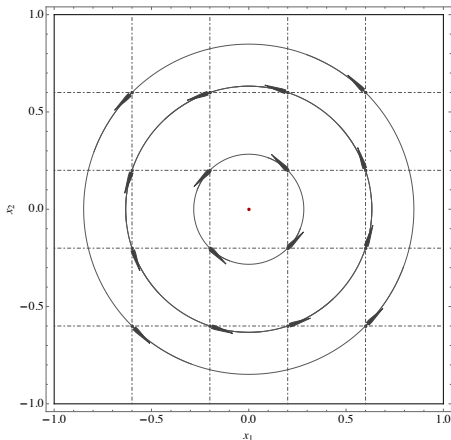


Figure: Hamiltonian flow of $f(x_1, x_2) = x_1 x_2$.

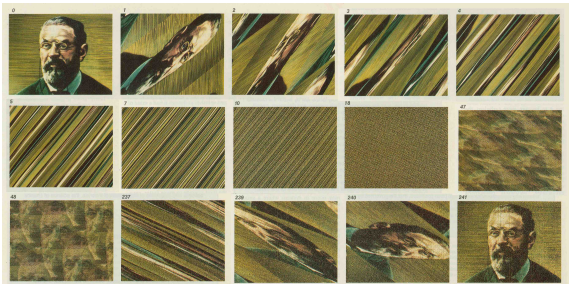


Poincaré recurrence

Cycles are an example of **recurrence**:

Definition (Poincaré, 1890's)

A dynamical system is *Poincaré recurrent* if almost all solution trajectories return *arbitrarily close* to their starting point *infinitely many times*.



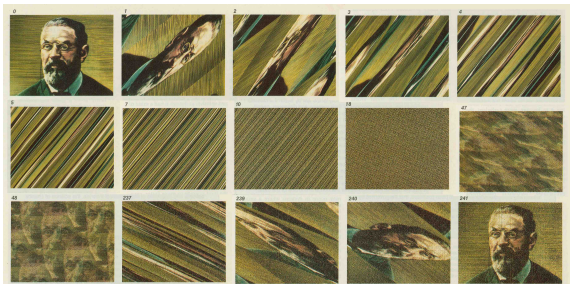


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Theorem (M, Papadimitriou, Piliouras, SODA 2018; bare-bones version)

(GD) is recurrent in all bilinear saddle-point problems with an interior solution.



OGD in games

OGD as a forward (Euler) scheme:

$$X^+ = X - \gamma V(X)$$



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Energy no longer a constant:

$$\frac{1}{2} \|X^+ - x^*\|^2 = \frac{1}{2} \|X - x^*\|^2 - \underbrace{\gamma \langle V(X), X - x^* \rangle}_{\text{from (GD)}} + \frac{1}{2} \underbrace{\gamma^2 \|V(X)\|^2}_{\text{discretization error}}$$

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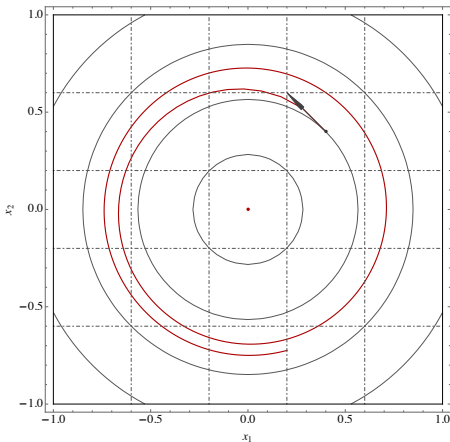
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OGD in games

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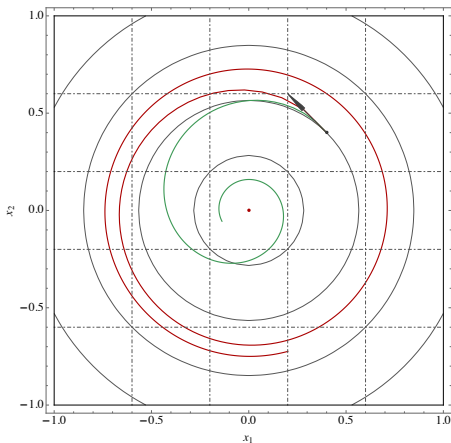




Time averages: a very different story

No-regret captures behavior of time-averaged process:

$$\bar{X}_t = \frac{1}{t} \sum_{s=1}^t X_s$$





Convergence to equilibrium

Behavior different under **strict** monotonicity:

$$\frac{1}{2} \|X_{t+1} - x^*\|^2 = \frac{1}{2} \|X_t - x^*\|^2 - \underbrace{\gamma_t \langle V(X_t), X_t - x^* \rangle}_{< 0 \text{ if } X_t \text{ not Nash}} + \frac{1}{2} \underbrace{\gamma_t^2 \|V(X_t)\|^2}_{\text{discretization error}}$$

Can the drift overcome the discretization error?



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Theorem (M & Zhou, MathProg 2019)

- ▶ **Assume:** game strictly monotone, $\sum_t \gamma_t = \infty$, $\sum_t \gamma_t^2 < \infty$
- ▶ **Then:** X_t converges to a Nash equilibrium from any initial condition



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In strictly monotone games, no-regret \rightsquigarrow Nash equilibrium



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Feedback

(OGD) requires gradient information, which may be difficult to come by:

- ▶ Other players' actions unknown
- ▶ Measurement errors
- ▶ Stochastic utilities (realized vs. expected gradients)
- ▶ ...

Imperfect gradient feedback:

$$\hat{V}_t = V(x_t) + U_t$$

with the following hypotheses:

$$\text{[H1] Zero-mean error: } \mathbb{E}[U_t \mid \mathcal{F}_{t-1}] = 0 \quad [\implies \mathbb{E}[\hat{V}_t \mid \mathcal{F}_{t-1}] = V(x_t)]$$

$$\text{[H2] Finite mean squared error: } \mathbb{E}[\|U_t\|_*^2 \mid \mathcal{F}_{t-1}] \leq \sigma^2 \quad [\implies \mathbb{E}[\|\hat{V}_t\|_*^2 \mid \mathcal{F}_{t-1}] \leq V^2]$$



Learning with imperfect gradients

Algorithm 1 Stochastic gradient descent

Require: step-size sequence $\gamma_t > 0$

- 1: choose $X \in \mathcal{X}$ # initialization
 - 2: **for** $t = 1, 2, \dots$ **do**
 - 3: oracle query at state X returns V # gradient feedback
 - 4: set $X \leftarrow \Pi(X + \gamma_t V)$ # new state
 - 5: **end for**
 - 6: **return** X
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Guarantees:

- ▶ $\mathbb{E}[\text{Reg}(T)] = \mathcal{O}(\sqrt{T})$ [folk]
- ▶ Strict monotonicity $\implies X_t$ converges to Nash (a.s.) [M & Zhou, 2019]



No gradient feedback whatsoever

In many cases, even stochastic gradients are out of reach:

- ▶ Multi-armed bandits (clinical trials, ...)
- ▶ Other players' actions unknown (auctions, ...)
- ▶ ...



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Possible fixes:

- ▶ Two-time-scale approach: fast samples, slow updates

[can be slow 😊]



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- ▶ Multiple-point estimates [needs synchronization ☹️]



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Possible fixes:

- ▶ Two-time-scale approach: fast samples, slow updates [can be slow ☹️]
- ▶ Multiple-point estimates [needs synchronization ☹️]
- ▶ **Simultaneous perturbation stochastic approximation** [Spall, 1997]



Simultaneous perturbation stochastic approximation

Estimate $u'(x)$ at target point $x \in \mathbb{R}$

$$u'(x) \approx \frac{u(x + \delta) - u(x - \delta)}{2\delta}$$



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$$\mathbb{E}[u(x + \delta z)z] = \frac{1}{2}u(x + \delta) - \frac{1}{2}u(x - \delta)$$

\implies Estimate $u'(x)$ up to $\mathcal{O}(\delta)$ by sampling u at $\hat{x} = x + \delta z$ and looking at $\frac{1}{\delta}u(\hat{x})z$



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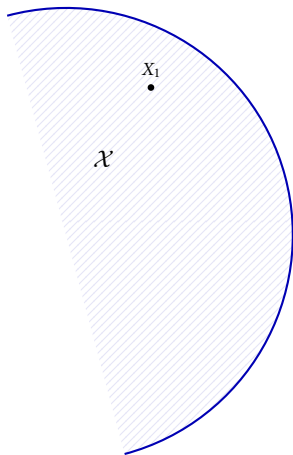
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Algorithm 2 Single-point estimator of ∇u at X

- 1: Draw z uniformly from \mathbb{S}^d
 - 2: Play $\hat{X} = X + \delta z$
 - 3: Get $\hat{u} = u(\hat{X})$
 - 4: Set $\hat{V} = \frac{d}{\delta} \hat{u} z$
-

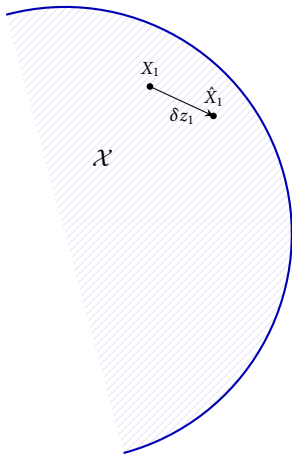


Learning with bandit feedback



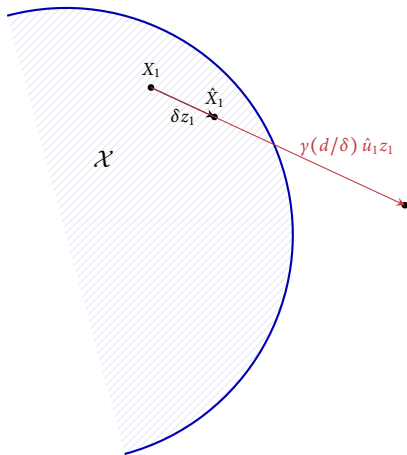


Learning with bandit feedback



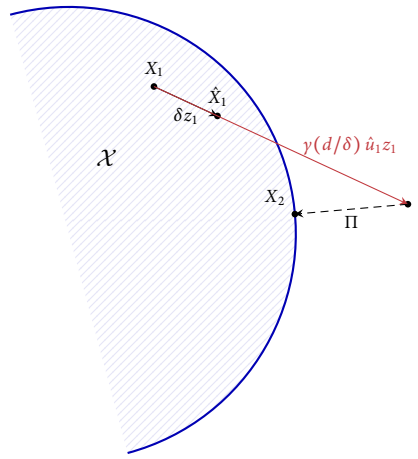


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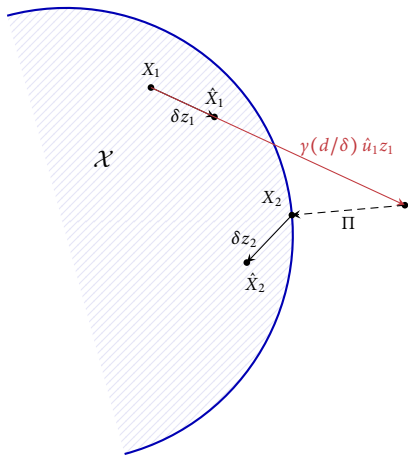


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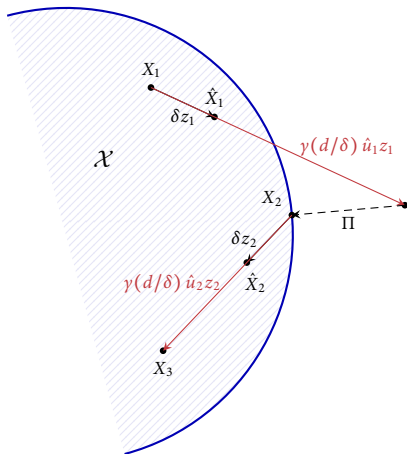


Learning with bandit feedback





Learning with bandit feedback





Bandit gradient descent

Algorithm 3 Multi-agent gradient ascent with bandit feedback

Require: step-size $\gamma_t > 0$, query radius $\delta_t > 0$, safety ball $\mathbb{B}_r(p) \subseteq \mathcal{X}$

- 1: choose $X \in \mathcal{X}$ # initialization
 - 2: **repeat** at each stage $t = 1, 2, \dots$
 - 3: draw Z uniformly from \mathbb{S}^d # perturbation direction
 - 4: set $W \leftarrow Z - r^{-1}(X - p)$ # feasibility adjustment
 - 5: play $\hat{X} \leftarrow X + \delta_t W$ # choose action
 - 6: receive $\hat{u} \leftarrow u(\hat{X})$ # get payoff
 - 7: set $\hat{V} \leftarrow (d/\delta_t)\hat{u} \cdot Z$ # estimate gradient
 - 8: update $X \leftarrow \Pi(X + \gamma_t \hat{V})$ # update pivot
 - 9: **until** end
-



Challenges

Key difficulty:

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- ▶ Stochastic approximation analysis requires **bounded variance**
- ▶ **Bias-variance dilemma**: accuracy vs. stability?



Convergence analysis

Must balance step-size γ_t against query radius δ_t :

- ▶ $\lim_{t \rightarrow \infty} \gamma_t = \lim_{t \rightarrow \infty} \delta_t = 0$ # vanishing noise and bias
- ▶ $\sum_{t=1}^{\infty} \gamma_t = \infty$ # the process doesn't stop
- ▶ $\sum_{t=1}^{\infty} \gamma_t^2 / \delta_t^2 < \infty$ # variance control
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Theorem (Bravo, Leslie & M, NIPS 2018)

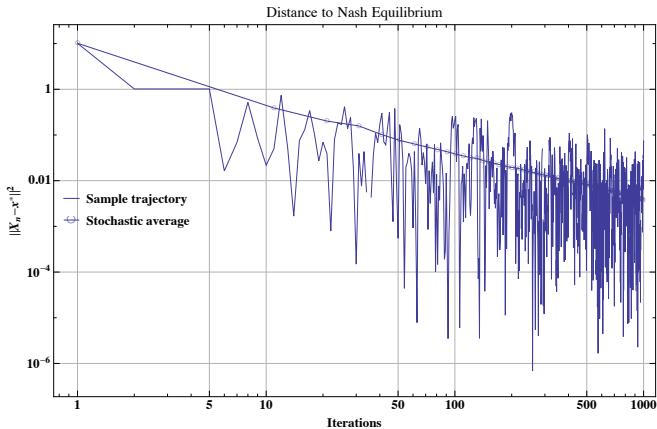
1. Under **strict monotonicity**, X_t converges to Nash equilibrium with probability 1.
2. Under **strong monotonicity** ($H(x) < -\beta I$), $\gamma_t \propto 1/t$, $\delta_t \propto 1/t^{1/3}$, we have:

$$\mathbb{E}[\|X_t - x^*\|^2] = \mathcal{O}(1/t^{1/3}).$$



Convergence rate

Speed of convergence in a repeated Kelly auction





Conclusions and perspectives

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Conclusions and perspectives

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Open questions

- ▶ Faster rates?
- ▶ Delayed payoff observations?
- ▶ Beyond monotonicity?
- ▶ ???

NetEcon 2019

The 14th Workshop on the Economics of Networks, Systems and Computation
Phoenix, Arizona, 28th June 2019

In conjunction with ACM EC 2019 & SIGMETRICS 2019

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